

# Universality for Timed Automata with Minimal Resources

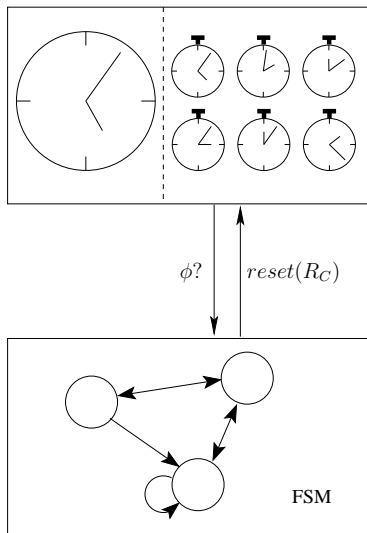
Sara E. Adams

in collaboration with  
Joël Ouaknine and James Worrell

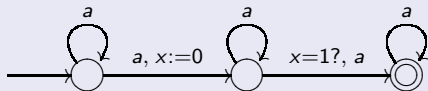
University of Oxford

FORMATS 2007

# Concept of a timed automaton



# An example timed automaton



## Why universality?

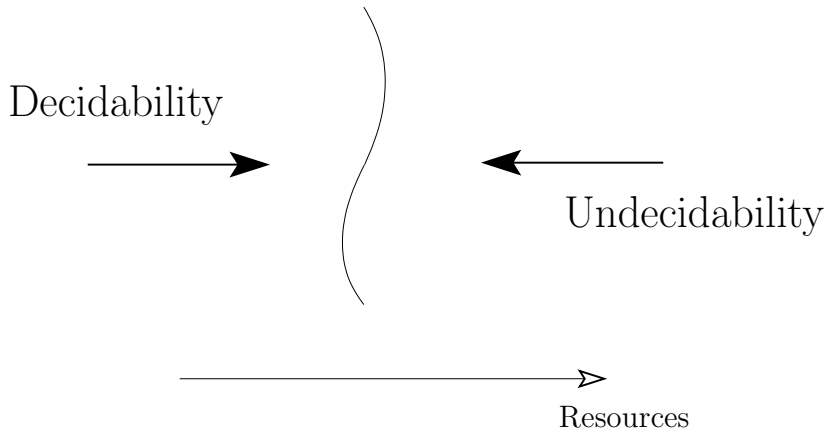
Special case of language inclusion

- All timed words  $\subseteq L \Rightarrow L$  universal
- Universality undecidable  $\Rightarrow$  language inclusion undecidable

## Why language inclusion?

Verification of real-time systems

- Essential role of language inclusion
- e.g. "Implementation  $\subseteq$  Specification"



# Universality Problem

*Does a given automaton accept every timed word?*

Alur and Dill, 1994

Universality is undecidable for timed automata with two clocks.

Ouaknine and Worrell, 2004

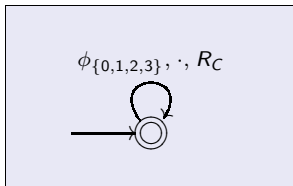
Universality is decidable for timed automata with one clock.

# Our contribution

## Adams, Ouaknine, Worrell

Universality is undecidable for timed automata with one state, one event and comparisons to

- weakly monotonic: 0 and 1 only.
- strongly monotonic: 1, 2, and 3 only.



## Obviously minimal

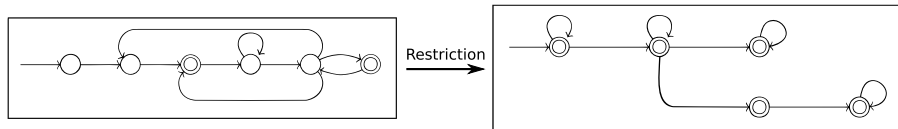
- essentially stateless
- essentially alphabetless

## Further observations

- comparisons to 0 only: decidable
- restrictions on number of clocks: decidable



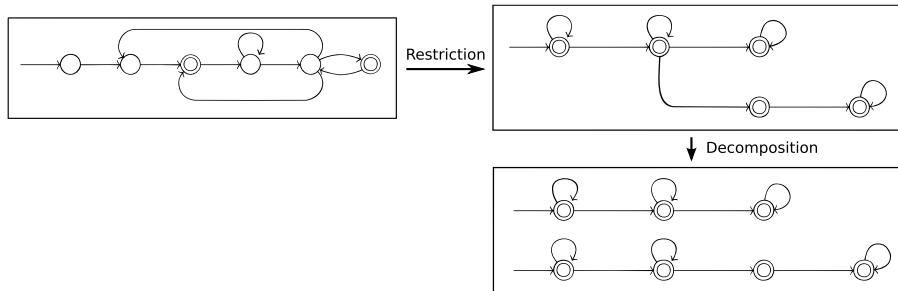
# Structure of the proof



## Basic steps

1. Universality for linear safety automata

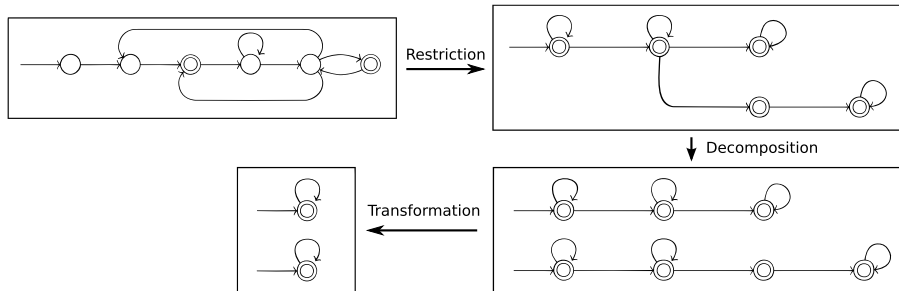
# Structure of the proof



## Basic steps

1. Universality for linear safety automata
2. Decomposition of linear safety automata

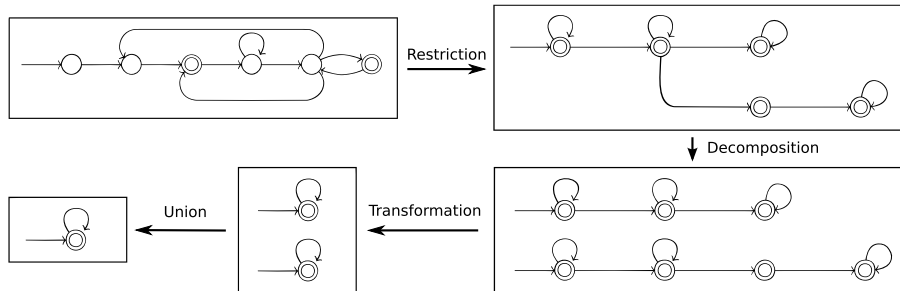
# Structure of the proof



## Basic steps

1. Universality for linear safety automata
2. Decomposition of linear safety automata
3. Transformation of decomposed timed automata

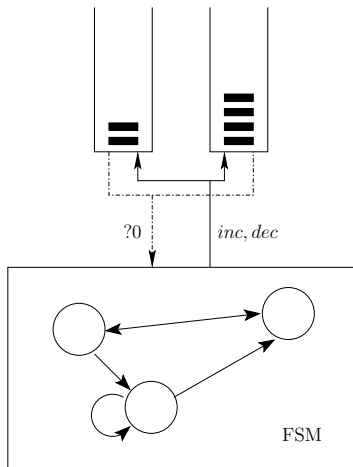
# Structure of the proof



## Basic steps

1. Universality for linear safety automata
2. Decomposition of linear safety automata
3. Transformation of decomposed timed automata
4. Union of transformed automata

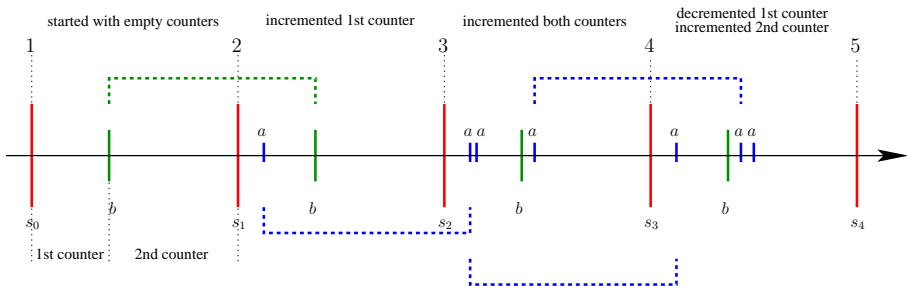
# Universality Proof: 2-counter machines



# Weakly monotonic time: 2-counter machine encoding

## Task

- Accept any inconsistent encoding of a halting computation
- Automaton universal  $\Leftrightarrow$  2-counter machine does not halt



# Weakly monotonic time: Converting to one symbol

## Simultaneous events

Encode alphabet symbols with simultaneous events

## Change of time model

before: strongly monotonic encoding

after: weakly monotonic encoding

# Weakly monotonic time: Converting to one state

## Clocks

For every state introduce a separate clock

## Rule

Use predicates to encode states:

- Reset state clock on state transition

## Essentials

- Linearity of automata
- On inconsistency: predicate that ensures acceptance



# Strongly monotonic time: 2-counter machine encoding

## Problem

Cannot use simultaneous events to encode alphabet

## Solution

Use only one symbol in the 2-counter machine encoding

## Price

Use 3 time units for each configuration:

- 1st time unit: encode the state
- 2nd time unit: encode the 1st counter
- 3rd time unit: encode the 2nd counter

Universality is undecidable for timed automata with

- a single state,
- a single event, and
- clock comparisons to
  - weakly monotonic: 0 and 1 only.
  - strongly monotonic: 1, 2, and 3 only.

# Open questions

## Weakly monotonic time

Is universality undecidable for timed automata with:  
one state, one event, and comparisons to 1 only?

## Strongly monotonic time

Is universality undecidable for timed automata with:  
one state, one event, and comparisons to one or two constants only?